

TeleTrust-interner Workshop

Bochum, 27./28.06.2013

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Impulsvortrag

Embedded Security for the Internet of Things **RUB**



TeleTrust Workshop
Zentrum für IT-Sicherheit
Bochum, 27.6. 2013

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Acknowledgement

- Benedikt Driessen
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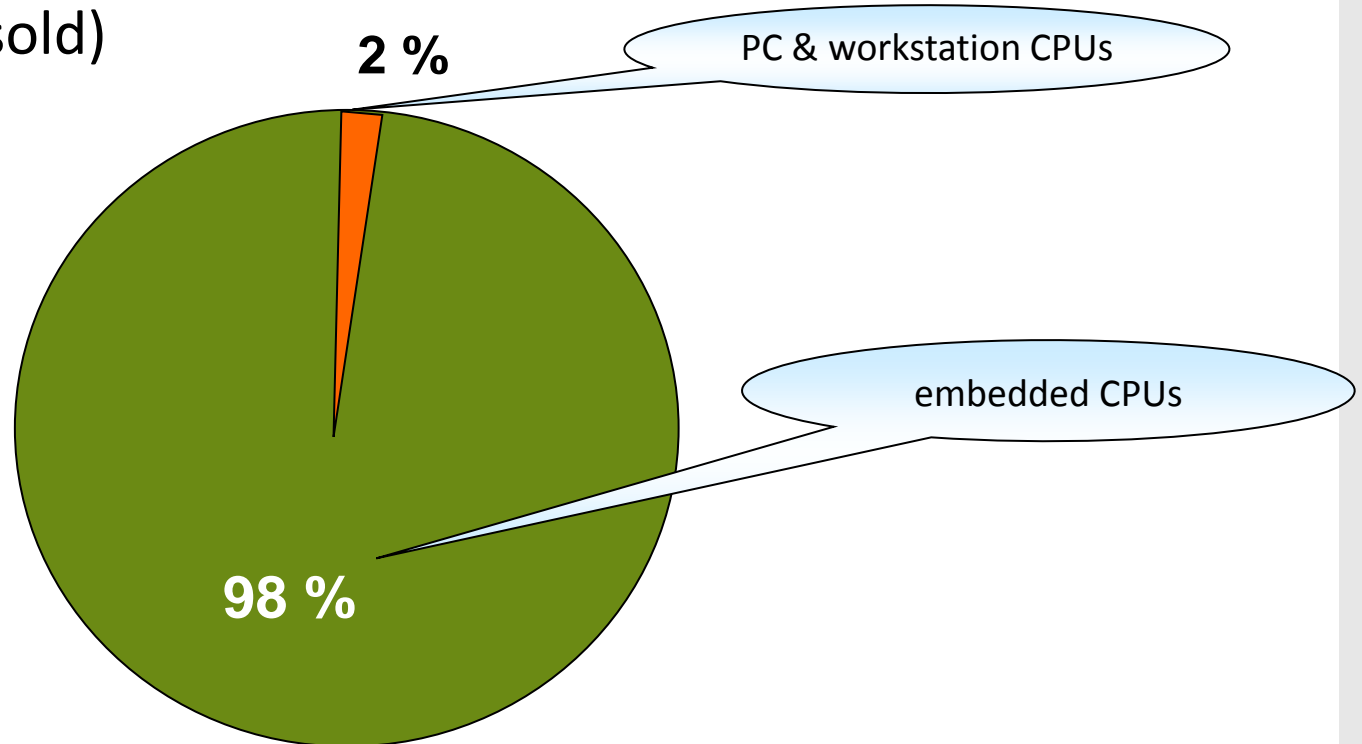
Agenda

- General Thoughts on Embedded Security
- Constructive: Bar Codes and SP Ciphers
- Destructive 1: Cell phones in the Desert
- Destructive 2: Routers and AES
- Auxiliary stuff

- **General Thoughts on Embedded Security**
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Who cares about *embedded systems*?

CPU market (units sold)



Q: But security ?

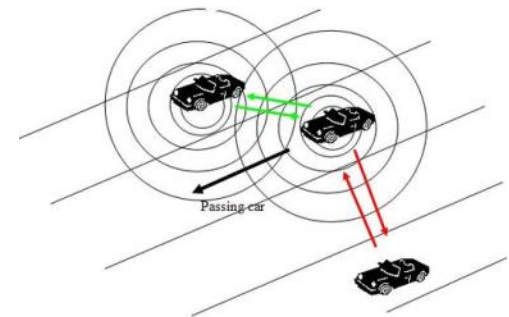
Embedded Security – Examples

Embedded DRM applications (iTunes, Kindle, ...)



Telemedicine

Privacy & security of car2car communication



Electronic IDs and e-health cards

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Lightweight Cryptography

1. “We need RFID security with less than 2000 gates”
Sanjay Sarma, AUTO-ID Labs, CHES 2002



2. Securing sensor networks, e.g., infrastructure

3. US\$ 3 trillions annually due to product piracy* (> US budget)

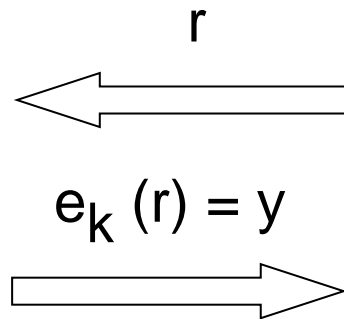


*Source: www.bascap.com

Needs authentication & identification

⇒ can both be fixed with standard cryptography

Strong Identification (symmetric crypto)



1. random challenge r
2. encrypted response y
3. verification
 $e_k(r) = y'$
 $y == y'$

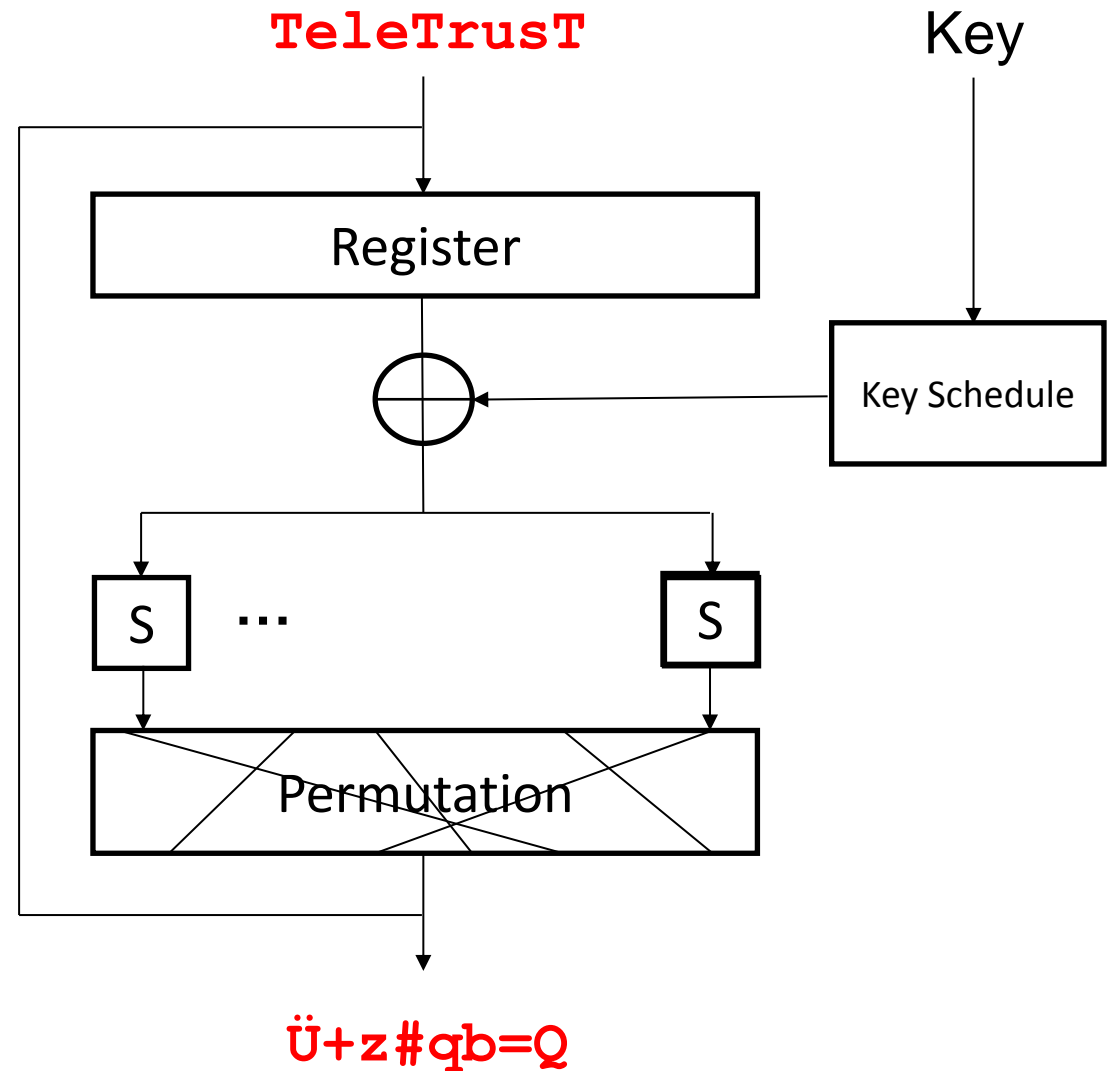
Challenge: Encryption function $e()$ at extremely low cost

→ almost all existing ciphers not optimized for cost ...

→ **Q: How cheap can we make cryptography?**

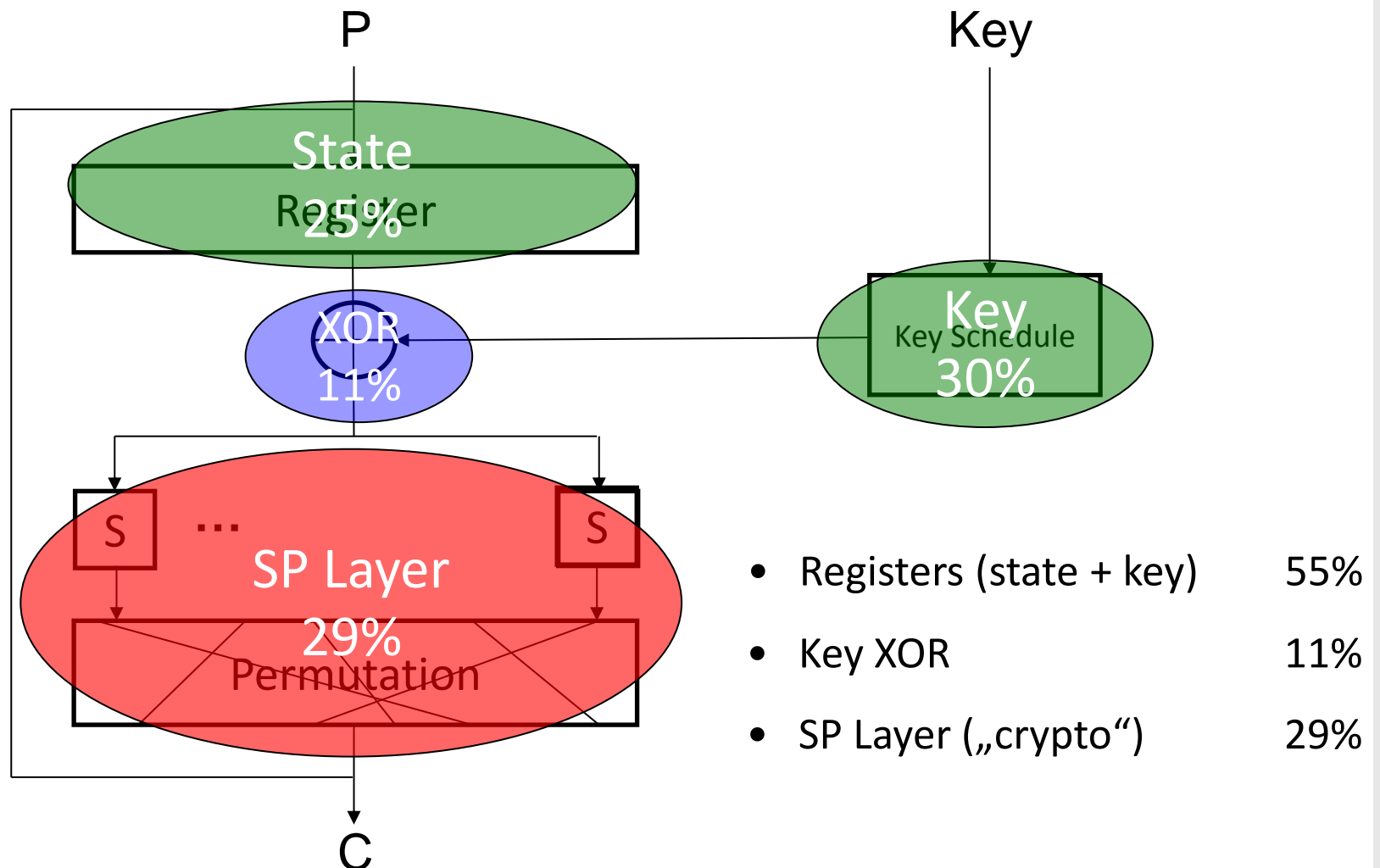
PRESENT – An aggressively cost-optimized block cipher for RFID

- 64 bit block, 80/128 bit key
- pure substitution-permutation network
- 4-4 bit Sbox
- 31 round (32 clks)
- secure against all known attacks
- joint work with Lars Knudsen, Gregor Leander, Matt Robshaw, ...

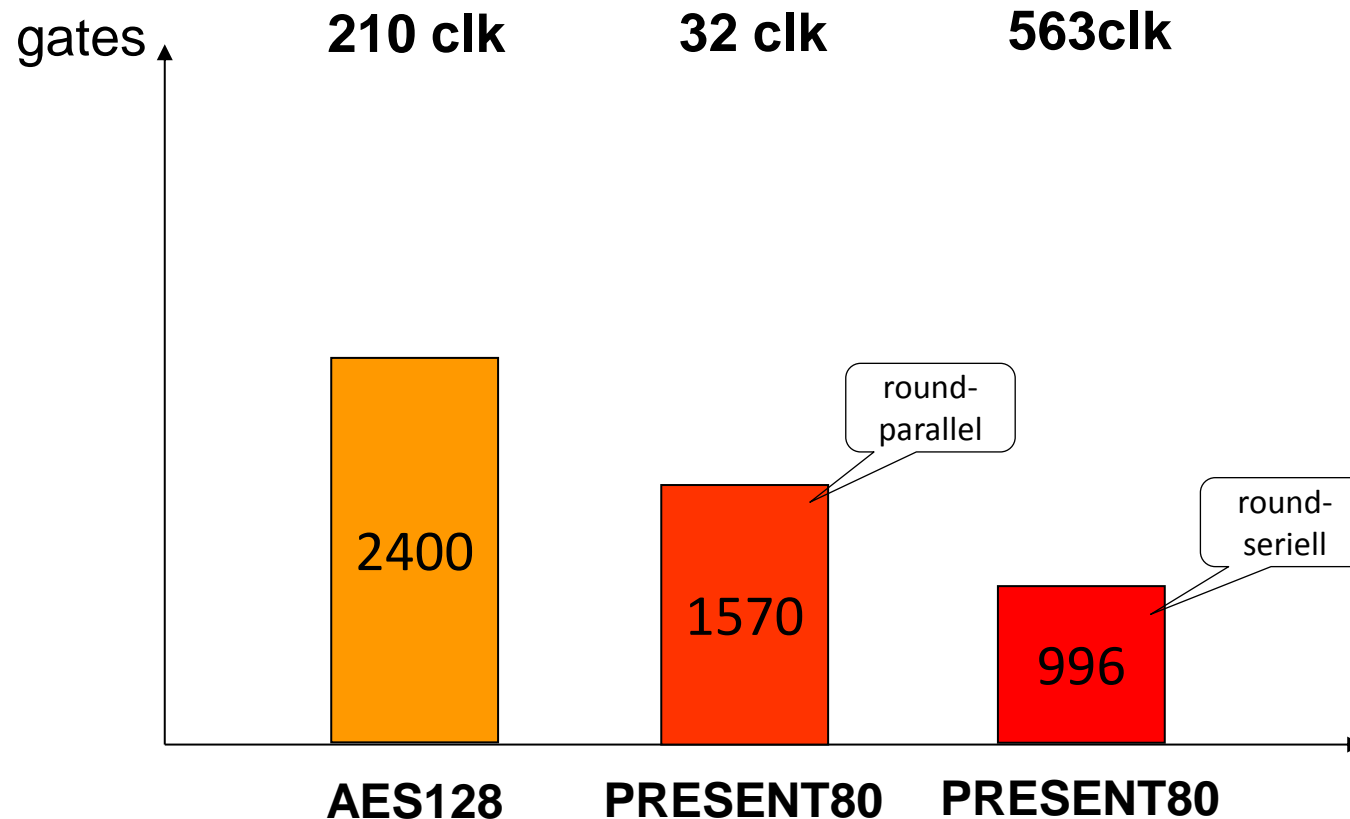


Resource use within PRESENT

Round-parallel implementation (1570ge)



Results – PRESENT



- $\approx 90\%$ less energy than small AES
- smallest secure cipher (20+ cryptanalytical publications)
- ISO standardized (2012)
- Serial implementation approaches theoretical complexity limit: almost all area is used for the 144 bit state (key + data path)
- Many related proposals: CLEFIA, Hight, KATAN, KTANTAN, Klein, mCrypton, Piccolo, ...

Further Reading

- Bogdanov, Knudsen, Leander, P, Poschmann, Robshaw, Seurin, Vikkelsoe: PRESENT: An Ultra-Lightweight Block Cipher. CHES 2007.
- Rolfes, Poschmann, Leander, P: Ultra-Lightweight Implementations for Smart Devices – Security for 1000 Gate Equivalents. CARDIS 2008.
- Borghoff et al.: PRINCE - A Low-Latency Block Cipher for Pervasive Computing Applications. ASIACRYPT 2012.

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Mobile Satellite Telephony

- Cellphone communication not available in many **remote places**

- crew on oil rig or ships on open sea
- airplanes
- many humanitarian missions
- research expeditions
- and many military applications ...



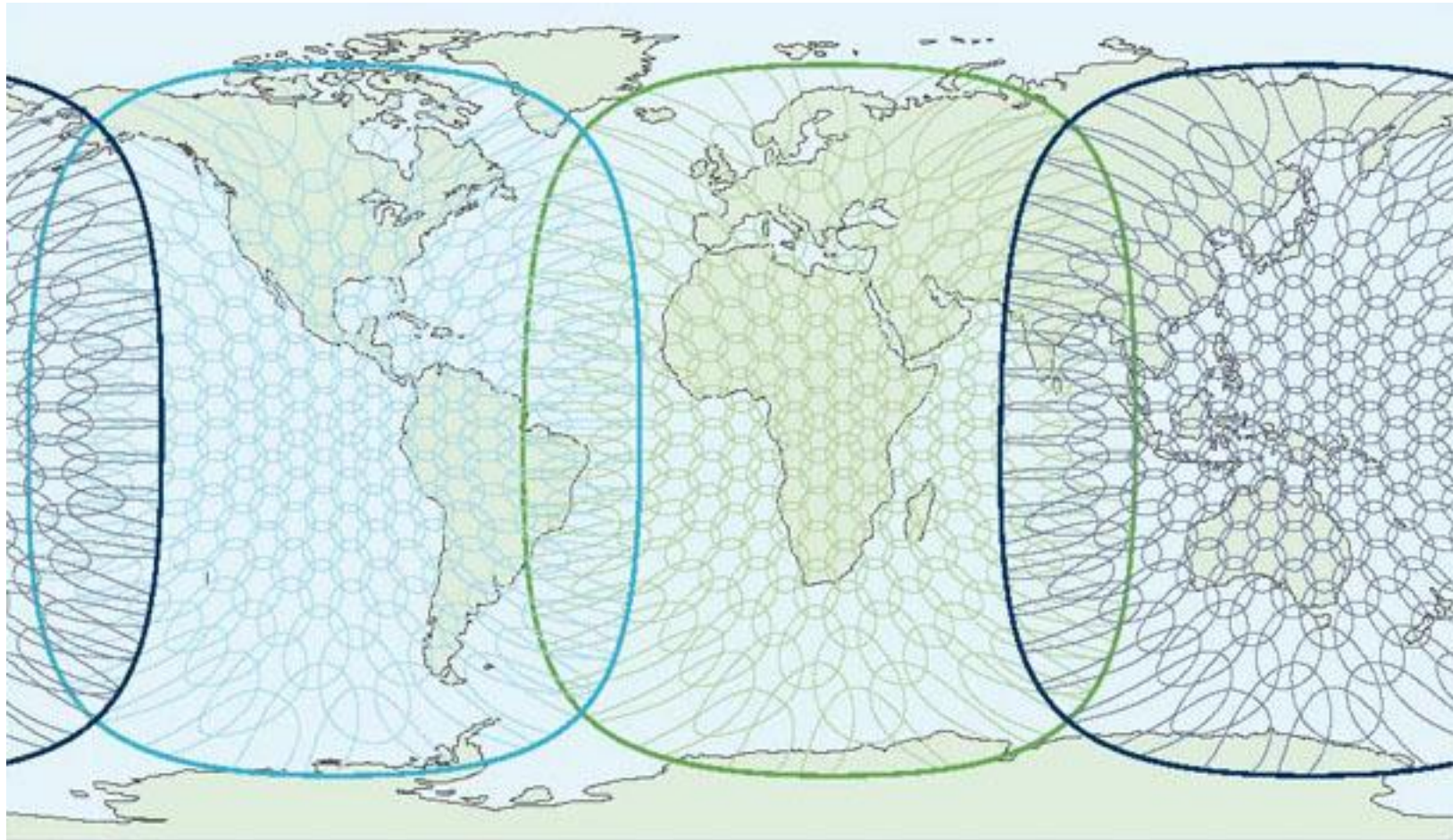
- **Satellite** telephony has been around since the 1990s
- **Direct** communication between phones and satellites

Satellite Phones

Many models



Inmarsat Spotbeam Coverage



Standards and Specifications

- Satellite phone systems standardized by **ETSI**
- 2 major standards coexist:
 - **GMR-1**: based on GSM, de-facto standard used by most providers
 - **GMR-2** (aka GMR-2+): based on GSM, used by Inmarsat (and ACeS)
- Specifications are **freely** available from ETSI
 - Both standards are very close to **GSM**
 - Cover signaling, encoding, etc.
 - ... **except the security parts of the standard**



Is Satphone Communication Secure?

Starting situation

- GSM algorithms have been (essentially) **broken**
- Q: Are the GMR algorithms **vulnerable** to similar attacks?

Research statement

Identify/extract the A5-GMR algorithms from satphones and perform cryptanalysis

Choosing a Target

- Several GMR-1 phones on the market
- Our victim: **Thuraya SO-2510** (for no specific reason)
- Firmware update publically available
- We didn't (have to) look at any other GMR-1 firmware
- Analysis was done **statically** only, we had no real phone at our disposal!



General Attack Procedure



1) Dump firmware image from firmware updater



2) Obtain information on the phone's actual hardware architecture



3) Dump DSP code from firmware image



4) Find cipher code

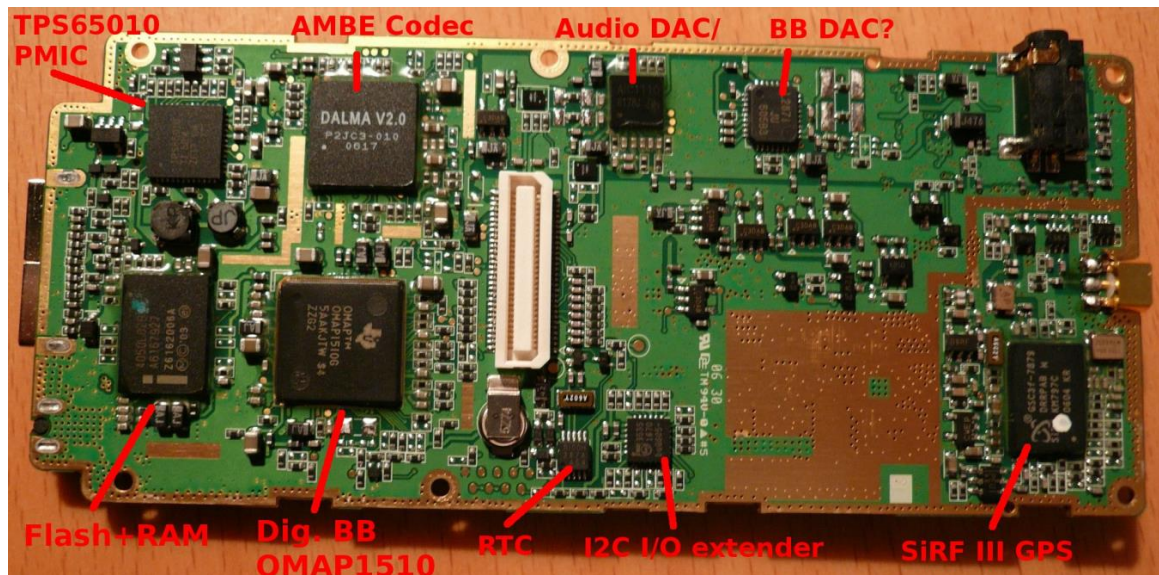


5) Perform cryptanalysis

Reverse engineering

Firmware and Hardware

- Firmware provided **unencrypted/unpacked** (some meta data needed to be stripped)
- Thuraya SO-2510 runs a TI **OMAP1510** platform (aka OMAP5910): ARM + TI C55X DSP
- fairly well documented platform
- OS is **VxWorks**



- Surprisingly, firmware image contains plenty of **assertion/log strings**
- Allows to deduce the name of some functions

```
ROM:0139B094 ; ===== S U B R O U T I N E =====
ROM:0139B094
ROM:0139B094
ROM:0139B094 vptr_slist_last ; CODE XREF: v_grm_dl_initialise_reassembly_buffer+1E4↓p
ROM:0139B094 var_14 = -0x14
ROM:0139B094
ROM:0139B094 STMFD SP!, {R4-R6,LR}
ROM:0139B098 SUB SP, SP, #4
ROM:0139B09C MOV R5, R0
ROM:0139B0A0 MOV R4, R1
ROM:0139B0A4 MOV R6, #0
ROM:0139B0A8 LDR R3, =aEnteringFunctionVptr__0 ; "Entering function vptr_slist_last()"
ROM:0139B0AC STR R3, [SP, #0x14+var_14]
ROM:0139B0B0 LDR R0, =0x201
ROM:0139B0B4 LDR R1, =a__CodeGrm_single_linkli ; "../code/grm_single_linklist.c"
ROM:0139B0B8 LDR R2, =0x16D
ROM:0139B0BC LDR R3, =aGrmS ; "[GRM]: %s"
ROM:0139B0C0 BL Log2
ROM:0139B0C4 CMP R5, R6
ROM:0139B0C8 BEQ loc_139B0E4
ROM:0139B0CC LDR R3, [R5, #8]
ROM:0139B0D0 STR R3, [R4]
ROM:0139B0D4 LDR R3, [R4]
ROM:0139B0D8 CMP R3, R6
ROM:0139B0DC LDRNE R3, [R4]
ROM:0139B0E0 LDRNE R6, [R3]
ROM:0139B0E4
ROM:0139B0E4 loc_139B0E4 ; CODE XREF: vptr_slist_last+34↑j
ROM:0139B0E4 LDR R3, =aExitingFunctionVptr__0 ; "Exiting function vptr_slist_last()"
ROM:0139B0E8 STR R3, [SP, #0x14+var_14]
ROM:0139B0EC LDR R0, =0x201
ROM:0139B0F0 LDR R1, =a__CodeGrm_single_linkli ; "../code/grm_single_linklist.c"
ROM:0139B0F4 LDR R2, =0x17F
ROM:0139B0F8 LDR R3, =aGrmS ; "[GRM]: %s"
ROM:0139B0FC BL Log2
ROM:0139B100 MOV R0, R6
ROM:0139B104 ADD SP, SP, #4
ROM:0139B108 LDMFD SP!, {R4-R6,PC}
ROM:0139B108 ; End of function vptr_slist_last
ROM:0139B108
ROM:0139B108 ; =====
```


Find Cipher Code

- Yields **240kb** of DSP code
- **TI C55x** assembler
 - Code hard to understand
 - Needs some initial training
- **No** strings, symbols, whatsoever
- How do we find the cipher code **conveniently**?
- A5-GSM cipher relies heavily on **linear feedback shift registers** (LFSRs)
- Typically uses many XOR and shift operations ...

```
ROM:19BA7      mov      AC0, T0
ROM:19BA9      mov      #0, AC0
ROM:19BAB      rptb     loc_19BE3
ROM:19BAE      add      *AR0+ << #16, AC0
ROM:19BB1      sftl     AC0, #-16, AC1
ROM:19BB4      xor      AC0 >> #14, AC1
ROM:19BB7      xor      AC0 >> #13, AC1
ROM:19BBA      xor      AC0 >> #10, AC1
ROM:19BBD      mov      AC1, AC2
ROM:19BBF      xor      AC0 >> #11, AC1
ROM:19BC2      xor      AC0 >> #15, AC2
ROM:19BC5      bfxpa    #5555h, AC1, T1
ROM:19BC9      bfxpa    #00000h, AC2, AC3
ROM:19BCD      or       T1, AC3
ROM:19BCF      mov      AC3, *AR1+
```

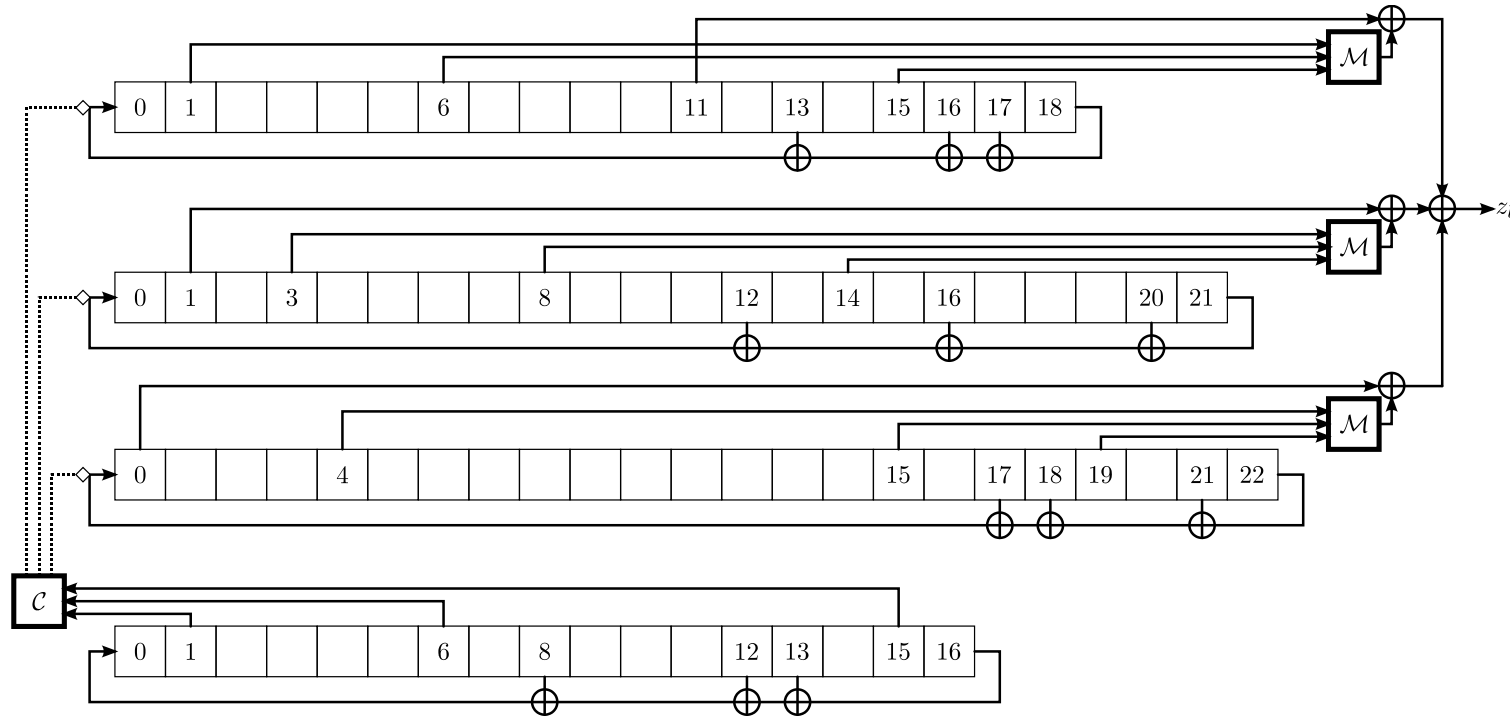
One identified function

```
sub_1CFC8:                                     ; CODE XREF: sub_1D0FC:loc_1D17C↓p  
                                              ; sub_1D24C+8B↓p  
      mov     db1(*abs16(#reg1)), AC1  
      sftl    AC1, #-1, AC2  
      mov     db1(*abs16(#reg1)), AC1  
      sftl    AC1, #-3, AC3  
      mov     db1(*abs16(#reg1)), AC1  
      xor     AC3, AC1  
      bfxtr   #0FFCh, AC1, AR1  
      mov     db1(*abs16(#reg1)), AC1  
      and     #1, AR1, AC3  
      xor     AC2, AC1  
      and     #1, AC1, AC1  
      xor     AC3, AC1  
      xor     AC0, AC1  
      sftl    AC1, #18, AC0  
      xor     AC2, AC0  
      mov     AC0, db1(*abs16(#reg1))  
      ret  
; End of function sub_1CFC8
```

- 4 such functions exist
- Each function does **exactly** one LFSR operation
- Reverse engineering of the 4 functions reveals the cipher...

The Cipher !

Looks familiar...



- A5-GMR-1 is basically „A5/2-GSM“
 - Feedback polynomials changed
 - Position of output taps changed
 - Initialization process changed slightly

Cryptanalysis

Ciphertext-Only Attack

- Ciphertext-only attack is possible
 - Based on ideas and [Barkan/Biham/Keller 2003]
- Adapt attack to GMR-1
 - Guess parts of R1, R2 and R3 to reduce variables and equations (increases the number of guesses...)

Results

- **Attack on voice channel possible with 16 frames + 2^{21} guesses**
- Experimental set-up cf. next slides

Attack Set-up with Software-Defined Radio

RUB



Further Reading

- Driessen, Hund, Willems, P, Holz: Don't Trust Satellite Phones: A Security Analysis of Two Satphone Standards.
IEEE Symposium on Security and Privacy 2012

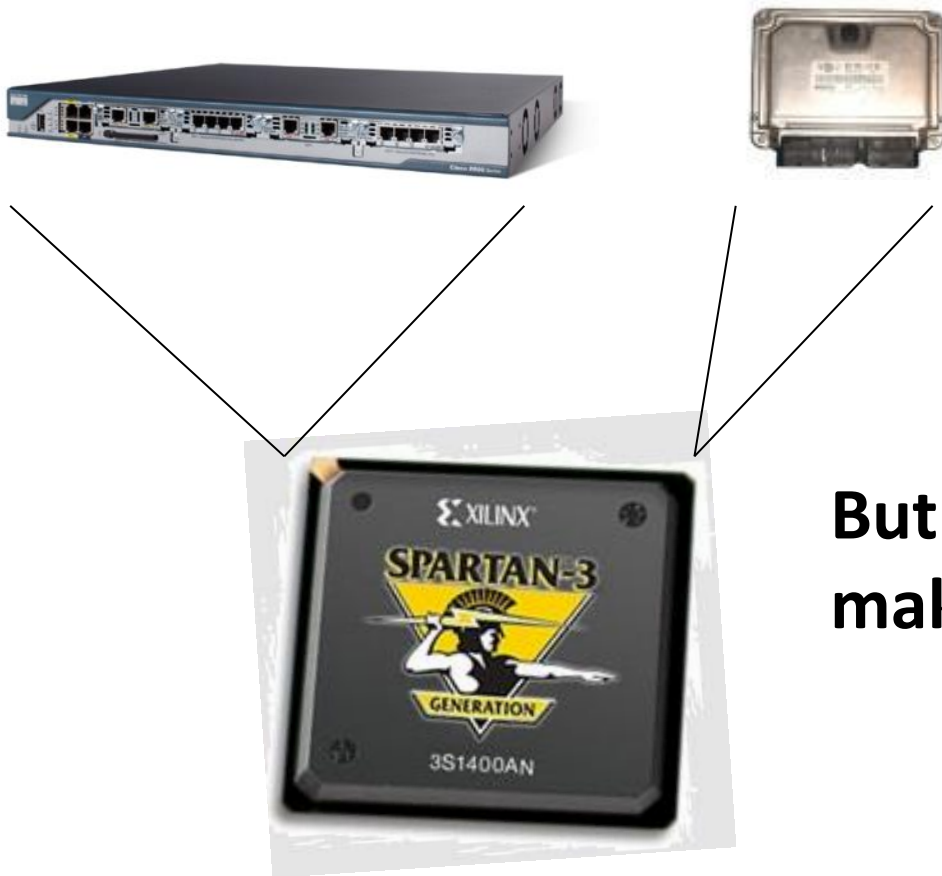
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FPGAs = Reconfigurable Hardware

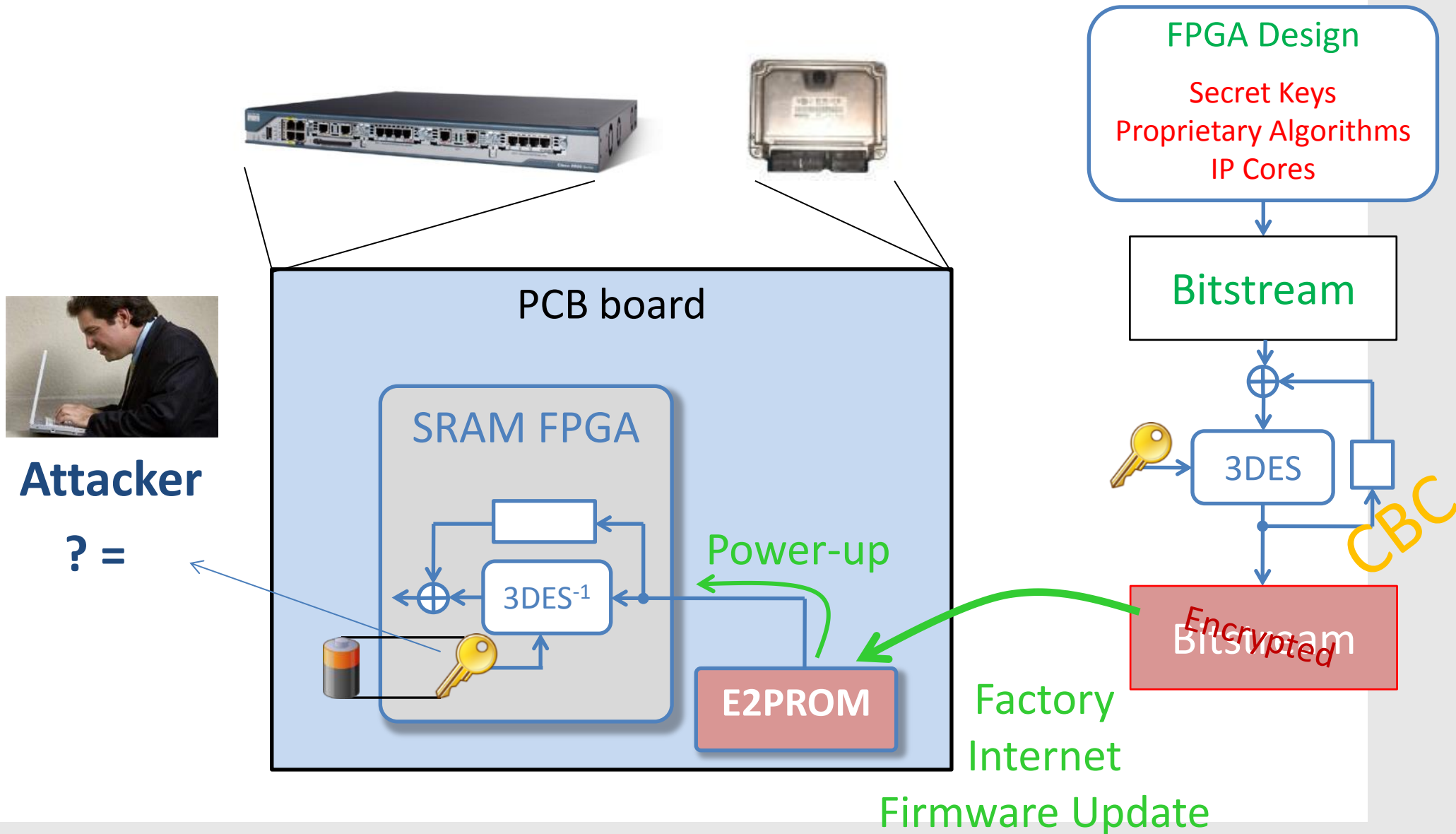
Widely used in

- routers
- consumer products
- automotive, machinery
- military

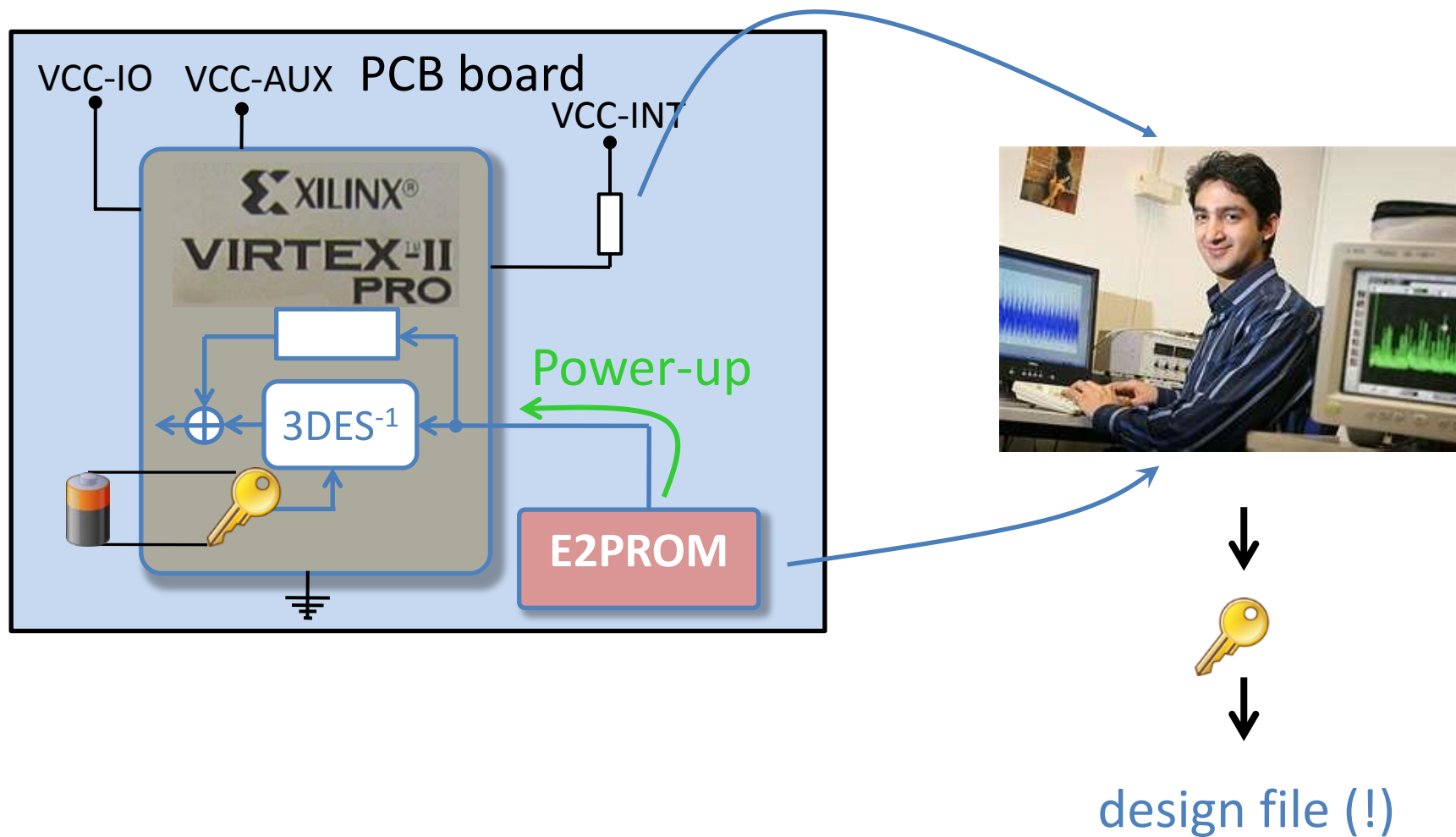


But: Copying the configuration files makes hardware counterfeiting easy!

Solution: Bitstream encryption



Let's try side-channel analysis

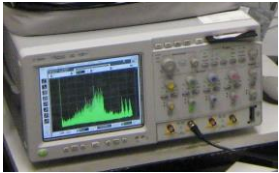


Side-Channel Attacks (1-slide version)



Analyze cipher

- Find a suited predictable intermediate value in the cipher



Measurements

- Measure the power consumption



Post Processing

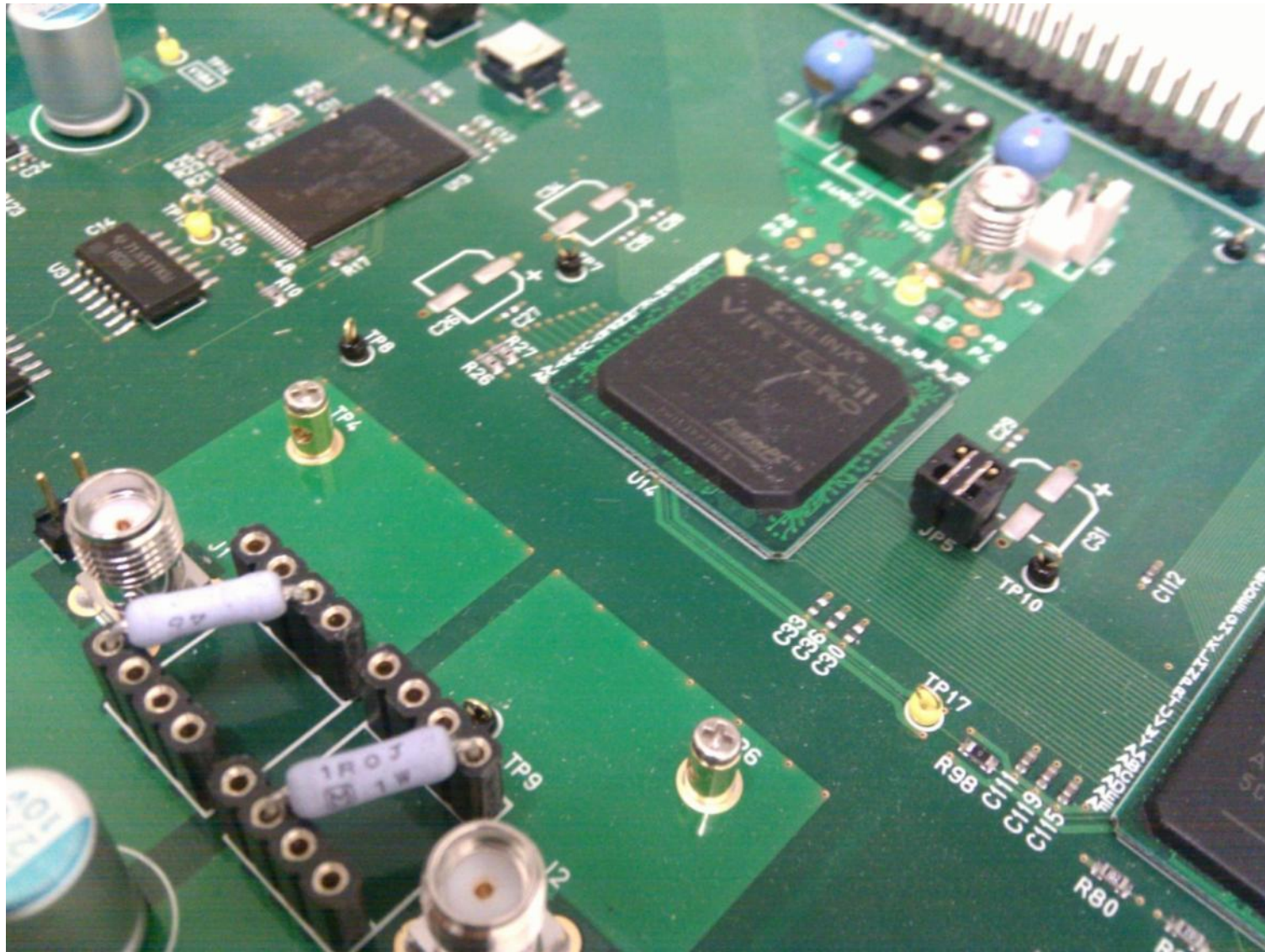
- Post-process acquired data



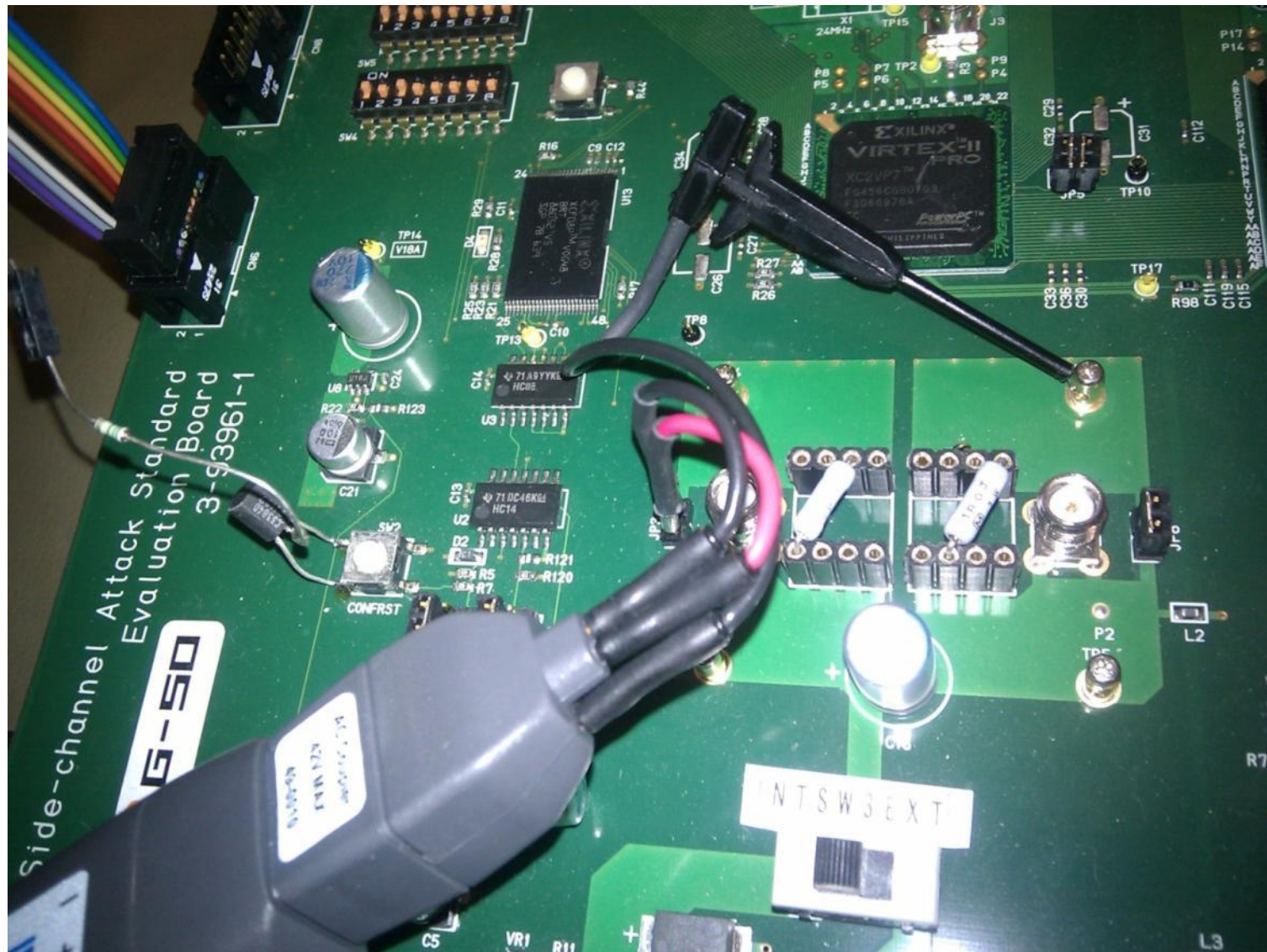
Key Recovery

- Perform the attack to recover the key

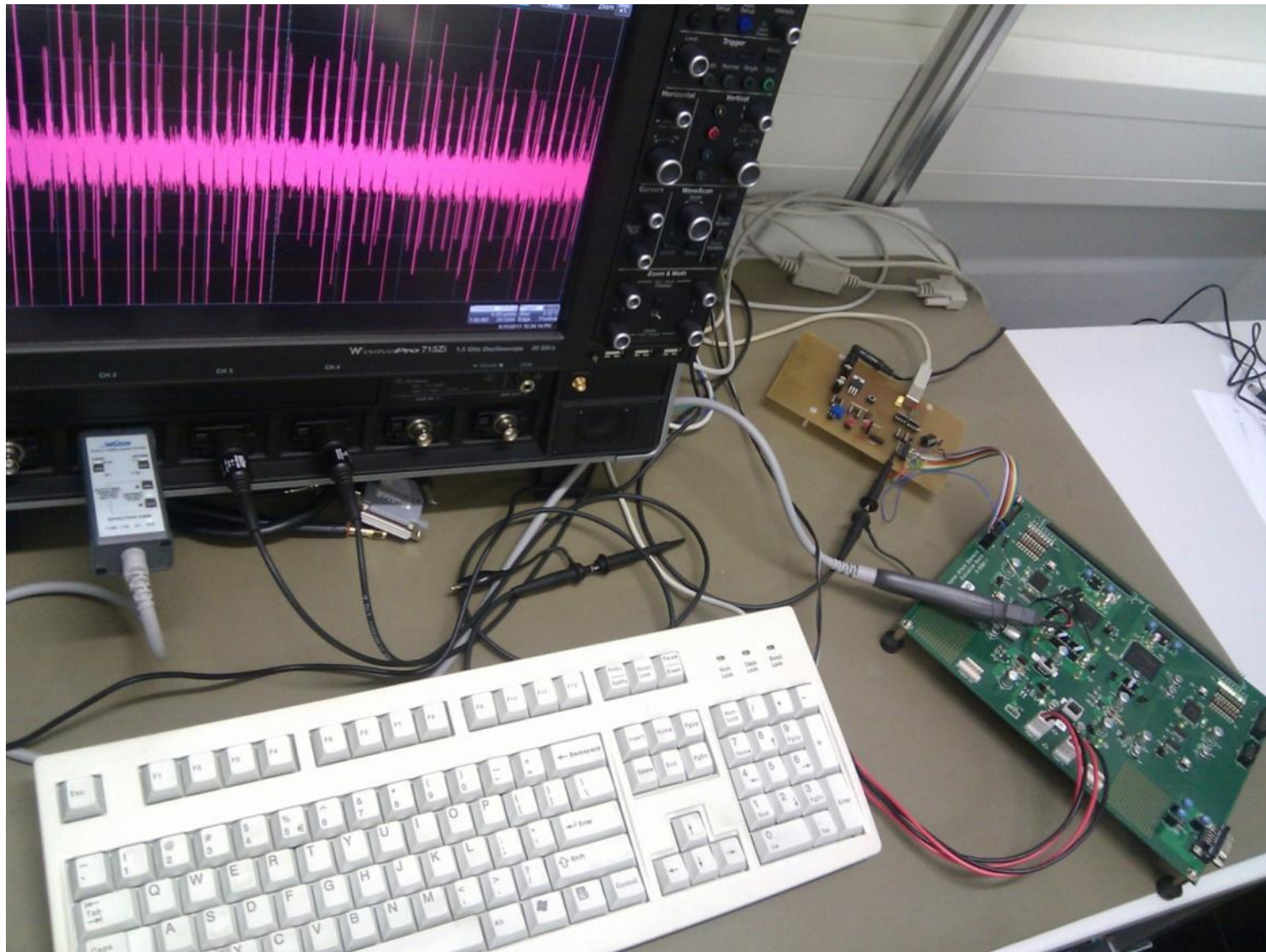
Our measurement set-up



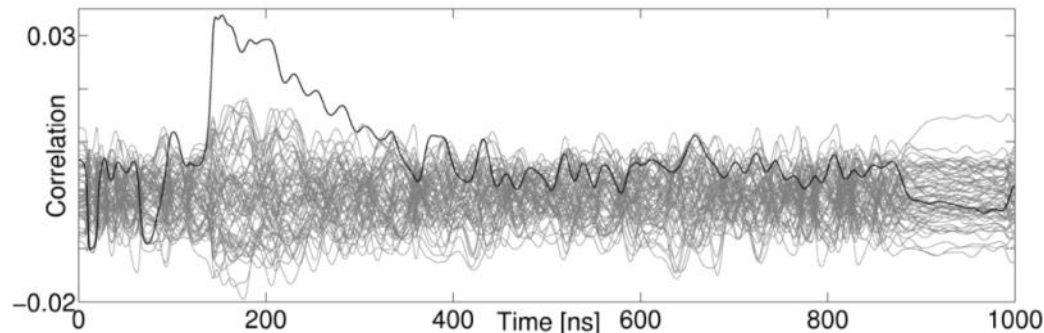
Our measurement set-up



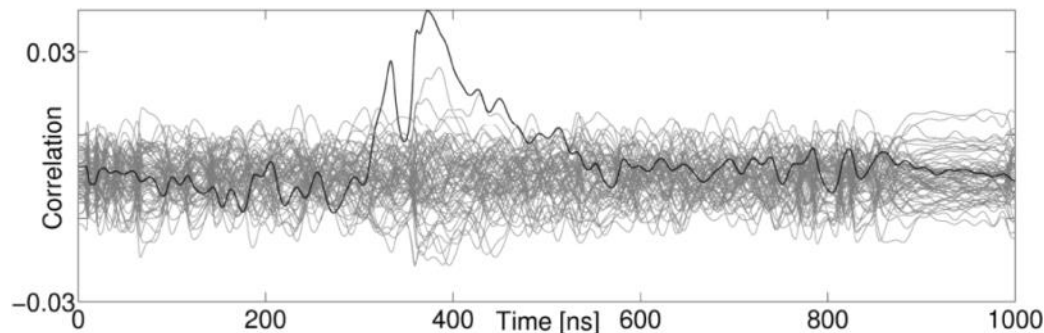
Signal acquisition



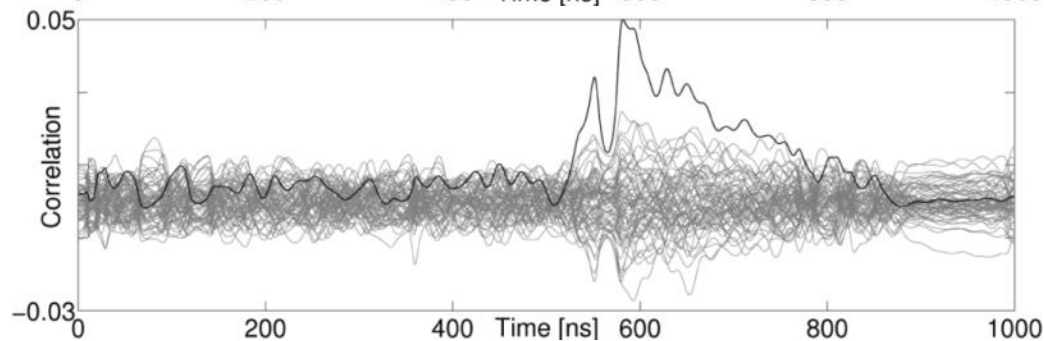
... 6 months later



key of 1st DES



key of 2nd DES



key of 3rd DES

Long story made short:

Decryption of “secret” designs is easy!

- Requires *single* power-up ($\approx 50,000$ traces)
- Complete 3DES key recovered with 2-3 min of computation
- Attack possible even though 3DES is only very small part of chip ($< 1\%$)
- Attack requires some experience, but
 - cheap equipment
 - easy to repeat

Implications

- Cloning of product
- Reverse engineering of design internals
- Alterations of design (chip tuning)
- Trojan hardware (i.e., malicious hardware functions)
- ...

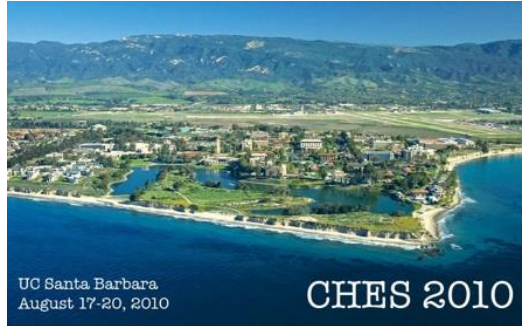
Further Reading

- Moradi, Barengi, Kasper, P: On the vulnerability of FPGA bitstream encryption against power analysis attacks: extracting keys from Xilinx Virtex-II FPGAs.
ACM CCS 2011
- Moradi, Oswald, P, Swierczynski: Side-channel attacks on the bitstream encryption mechanism of Altera Stratix II: facilitating black-box analysis using software reverse-engineering.
FPGA 2013

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Related Workshops



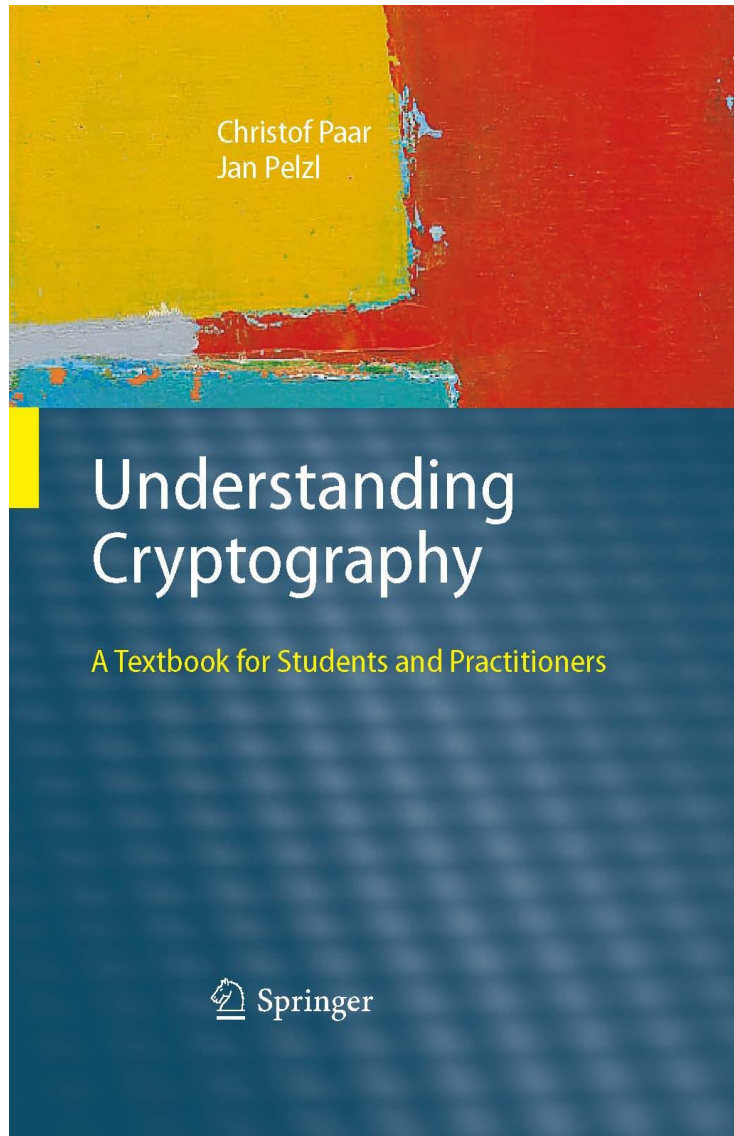
CHES – Cryptographic Hardware & Embedded Systems
August 2013, Santa Barbara, CA, USA

RFIDsec 2013
July 2013, Graz, Austria



escar USA – Embedded Security in Cars
November, Frankfurt, Germany

yet another textbook on cryptography (but this one targets engineers)



www.crypto-textbook.com

- includes videos of 2-semester course (in English)
- complete set of slides
- many further resources